CS 331, Fall 2024
Lecture 11 (10/7)

Today: - Graph search
- BFS
- DFS

Graph search (Part V, Section 1)

We have already seen some graph algos.

- SSSP on DAGs
- APSP (Floyd-Warshall)
- MST (Kruskal)

Basic theme:

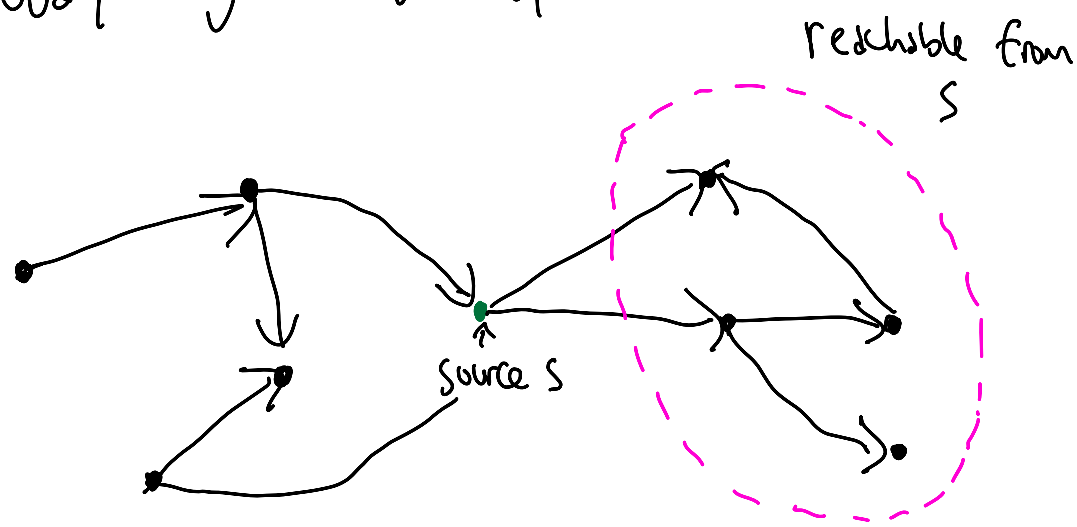
graph structure + algos bag
(lots to come!) (recursion, DP, greedy, data structures)

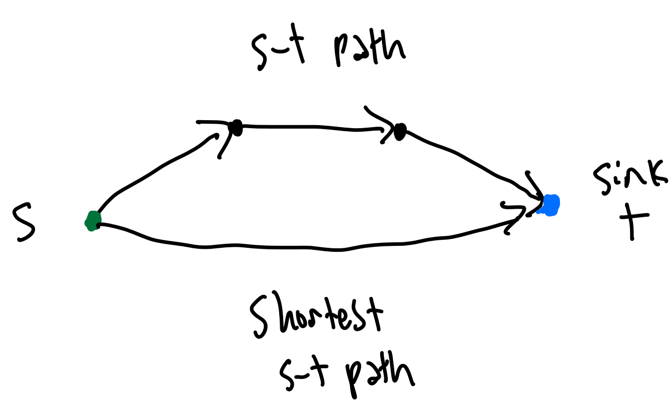
Example

MST: "greedy stays ahead"

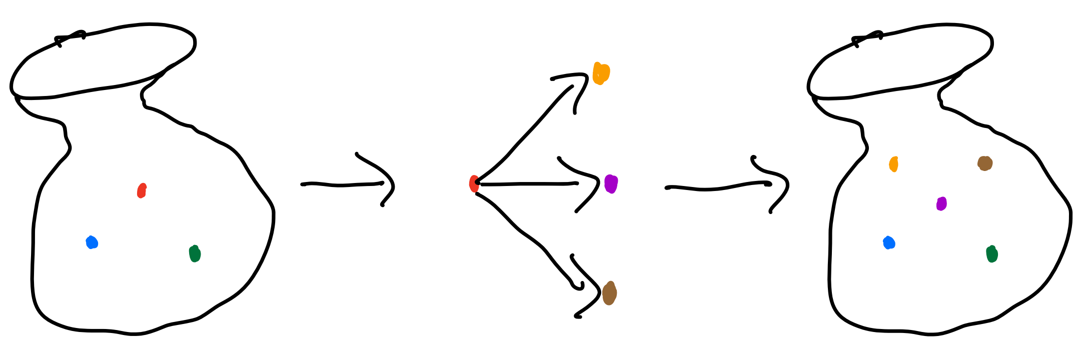
- exchange lemma: k CC's
  n-k edges
- data structure for maintaining cc's

Today: graph search





```
GraphSearch(G=(V,E), s):
  S ← {s}
  R ← [False for v ∈ V]
  While S ≠ ∅:
      v ← any element of S
      S ← S \ v
      If R[v] == False:
          R[v] ← True
          For (v,u) ∈ E: S ← S ∪ {u}
  Return R
```



Claim: GraphSearch solves reachability.

Proof: Every vertex set to True ≤ 1 time.

$R[v] = \text{True}$
$\Downarrow$
v reachable

Induct on when set to True.

Base case: s first.

Induct: If $R[v] = \text{True}$,

v was added b/c $R[u] = \text{True}$.

By assumption, u reachable $\Rightarrow$ so is v.

v reachable
$\Downarrow$
$R[v] = \text{True}$

Induct on shortest path distance.

If 0: $R[s] = \text{True}$.

If k: let path be $s \rightsquigarrow u \rightarrow v$

$R[u] = \text{True}$

Then, v added to S. Will become True

Breadth-First Search (Part V, Section 2.1)

How to implement Graph Search?

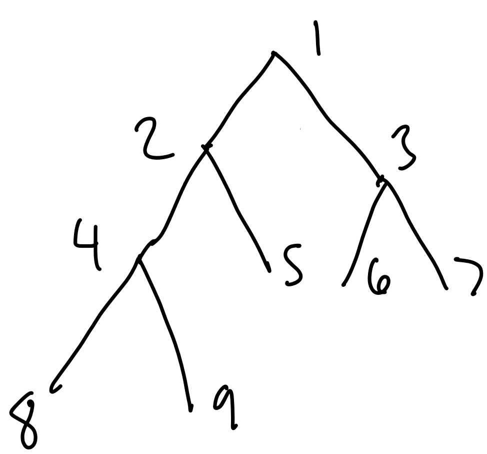
Need <u>data structure</u> for S.

- Insert an element
- Remove an element

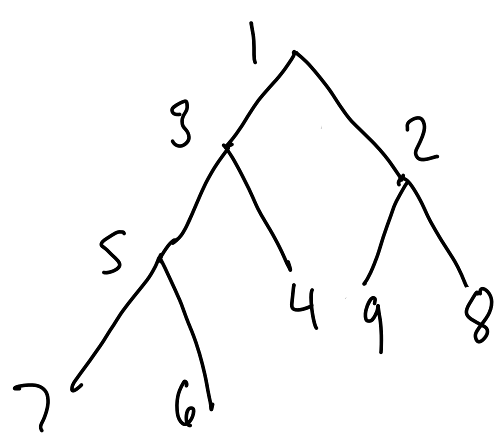
Idea: Linked List does both in $O(1)$ time.

Queue = BFS Push( ):

Stack = DFS Push( ):



BFS



DFS

Connected components

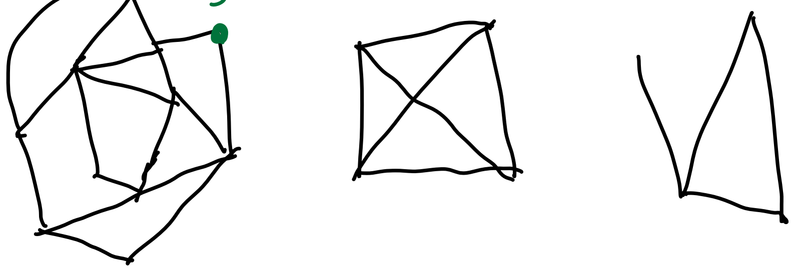
Say $s \sim t$ if connected in undirected graph

Then $\sim$ is equivalence relation

- reflexive 
- symmetric 
- transitive 

Partition into <u>connected components</u> (CCs)

s



reachable from s

Runtime of BFS: — every edge used ≤ 2x.

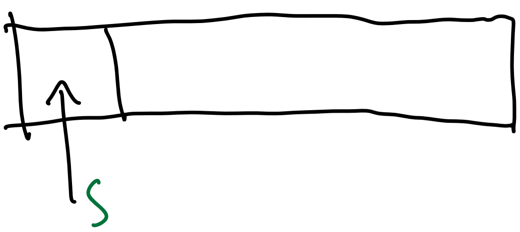
— $2 m_{C_s}$ total vertices

— m_{C_s} = # edges in C_s : CC of s.

$O(m_{C_s})$

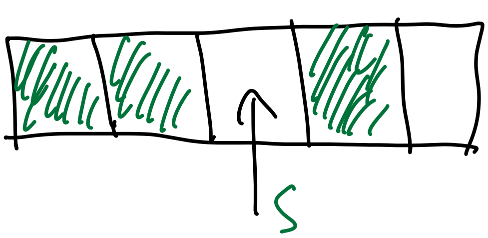
Using Queue (or Stack)

CC algo:



Vertex list: total move $O(n)$

GraphSearch(s) … $O(m_{C_s})$



$$O(n) + \underbrace{\sum O(m_{C_s})}_{O(m)}$$

Unweighted SSSP

Let $p(v)$ be the parent of v:

v added for the first time

due to the edge $(p(v), v)$

Claim: $\delta(s, v) = 1 + \delta(s, p(v))$

With Claim, simple mods compute SSSP!

- For $(u, v) \in E$: $S.\text{Push}((u, v))$

include parent info

- If $R[v] == \text{False}$:

$D[v] = D[u] + 1$

memoized

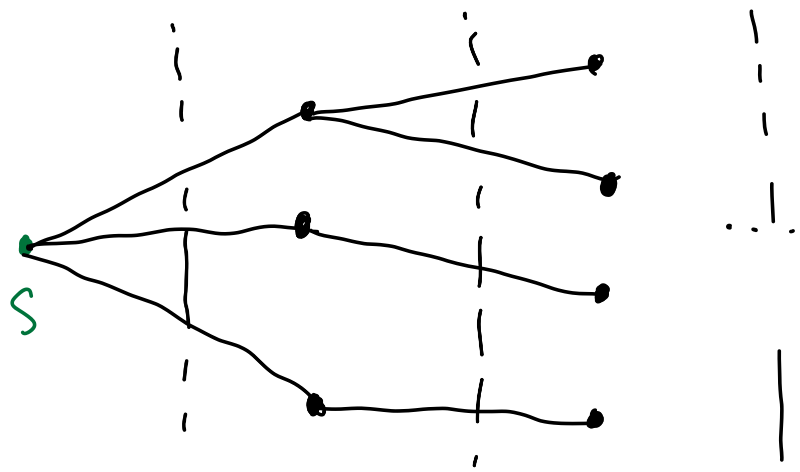
Proof of claim: Let $R_0 = \{s\}$

$R_1 = \{d(s, \cdot) = 1\}$

$R_2 = \{d(s, \cdot) = 2\}$

$\vdots$

Claim is that they form "frontiers" in S:



$$\underbrace{R_0 \quad \dots \quad R_1 \quad \dots \quad R_2 \quad \dots}_{\text{Queue, restricted to first visits}}$$

Base case: R_0 ✓

Induct: $R_0 \; R_1 \; R_2 \dots \mid R_i$

current queue

let $v \in R_i$ be dequeued

All (v, u) have $d(s, u) \leq i + 1$

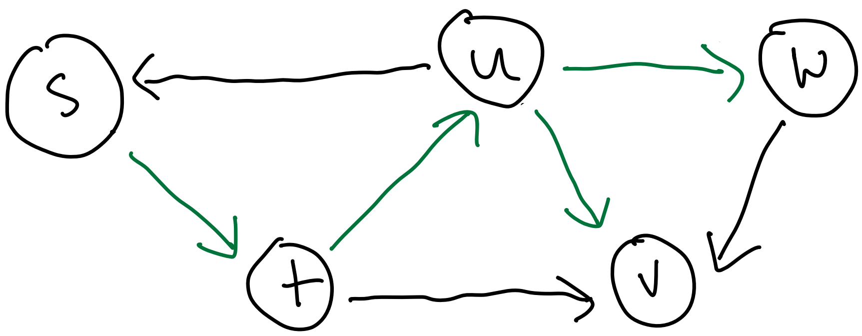
If $d(s, u) \leq i$ then reached (induction). ✓

Depth-First Search (Part V, Section 2.2)

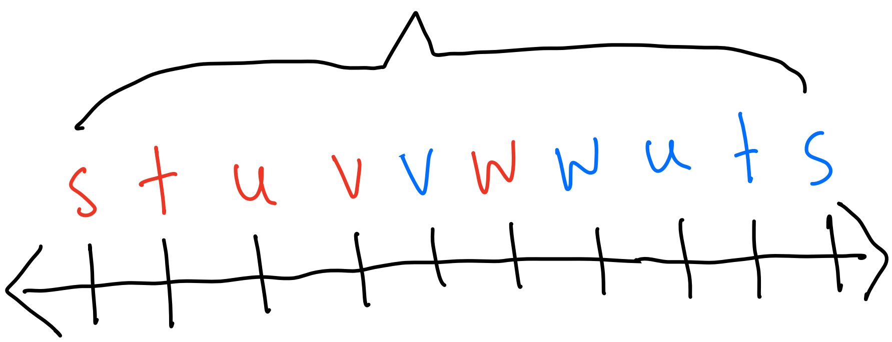
Each vertex has:

- start time (enter the stack)
- end time (leave the stack)
  - when all children done executing
  - implementable at no overhead (see notes)

Example



Stack: Duration of s on the stack



Preorder: s t u v w

Postorder: v w u t s

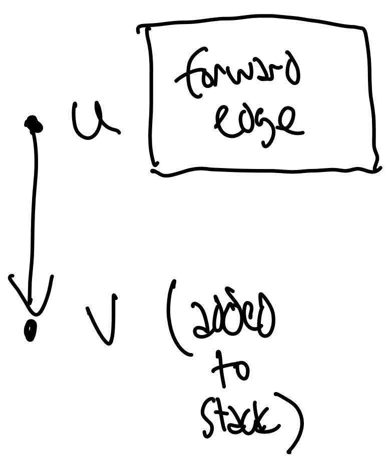
Key Claim:
If input is DAG, then postorder reversed = topological order!

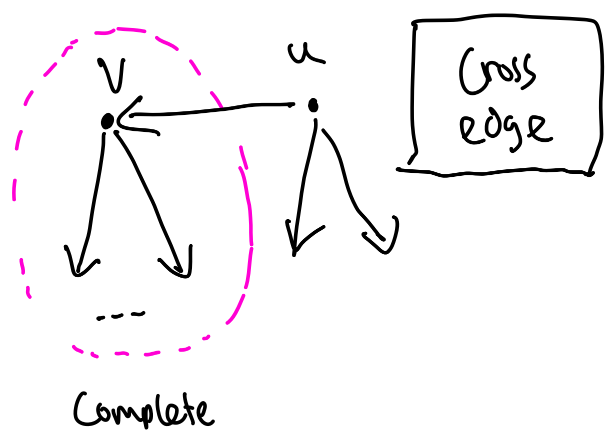
Lemma: Suppose $u \to v$,
u, v reachable from s.

If $u.\text{finish} < v.\text{finish}$, $\exists$ cycle

Proof: three cases when u reached

1. • v not visited ($u.\text{start} < v.\text{start}$)
2. • v finished ($v.\text{finish} < u.\text{start}$)
3. • v currently active ($v.\text{start} < u.\text{start} < v.\text{finish}$)

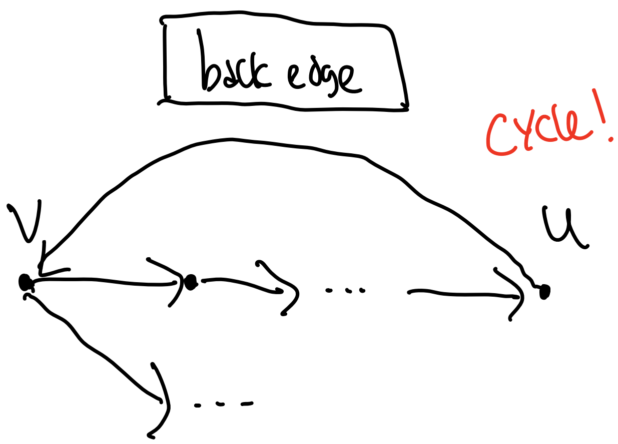




①: $v.\text{finish} < u.\text{finish}$

② same

(3) u belongs to v's recursive subtree



Only case with u.finish < v.finish $\blacksquare$

Punchline: suppose $u \to v$

u.finish < v.finish (top order failed?)

No! It's not a DAG by lemma.

If DAG: run DFS, reverse postorder

If not DAG: just check the edges!

$O(m+n)$ time.